

$$
\begin{aligned}
45_{10}=32+8+4+1 & =2^{5}+\underset{0}{0}+2^{3}+2^{2}+0+2^{0} \\
& =1
\end{aligned} 1 \quad 1 \quad 0 \quad 12
$$

10 is placed in the $2^{1}$ and $2^{4}$ positions, since all positions
for Another example is the following:

$$
\begin{aligned}
76_{10}=64+8+4 & =2^{6}+\begin{array}{l}
0 \\
0
\end{array}+\frac{2^{3}}{0}+2^{2}+0+0 \\
& =1
\end{aligned} 0 \quad 0 \quad 1 \quad 1 \quad 0 \quad 0
$$

- Mutltiply the given fractional decimal no by the base (radix) $r$
- Record the carry generated in this multiplication as MSD
- Multiply only the fractionsl number of the product in step 2 by the base and record the carry as the next bit of MSD
- Repeat steps 2 and 3 upto the end. The last carry will represent the LSD of equivalent binary no.



## Binary

- Values are 0,1 and are called bits
- Leftmost digit is MSB (Most Significant Bit) and Rightmost is LSB (Least Significant Bit)
- Applications: Mostly used in Digital Systems like Computers.

| Name | Size (bits) | Example |
| :---: | :---: | ---: | ---: |
| Bit | 1 | 1 |
| Nibble | 4 | 0101 |
| Byte | 8 | 00000101 |
| Word | 16 | 0000000000000101 |
| Double Word | 32 | 00000000000000000000000000000101 |

- Write down the number
- Write down the weights for different positions
- Multiply each digit in the given number with the corresponding weight to obtain product numbers
- Add all the product numbers to get the decimal equivalent.

$$
\begin{array}{ccccc}
1 & 1 & 0 & 1 & 1_{2} \\
2^{4}+ & 2^{3}+0 & 0+2^{1}+2^{0} & =16+8+2+1 \\
& =27_{10}
\end{array}
$$

Let's try another example with a greater number of bits:

$$
\begin{gathered}
1 \\
2^{7}+0+2^{5}+2^{4}+ \\
0
\end{gathered}+\begin{gathered}
0 \\
0
\end{gathered} 2^{2}+0+2^{0}=181_{10}=
$$

Convert $37_{10}$ to binary. Try to do it on your own before you look at the solution
Solution

$$
\begin{array}{cc}
\frac{37}{2}=18.5 \longrightarrow \text { remainder of } 1 \text { (LSB) } \\
\frac{18}{2}=9.0 \longrightarrow & 0 \\
\frac{9}{2}=4.5 \longrightarrow & 1 \\
\frac{4}{2}=2.0 \longrightarrow & 0 \\
\frac{2}{2}=1.0 \longrightarrow & 0 \\
\frac{1}{2}=0.5 \longrightarrow & 1(\mathrm{MSB})
\end{array}
$$

Thus, $37^{10}=\mathbf{1 0 0 1 0 1}_{2}$.

| S.NO | RGPV QUESTIONS | Year | Marks |
| :--- | :--- | :--- | :--- |
| Q.1 | Convert (27)10 and, Convert (40.25)10 to ()2, ()8, ()16 | Dec 2003 | 10 |
| Q. 2 | Convert (350.25)10 to ()2, ()8, ()16 | Dec 2005 | 6 |
| Q.3 | Convert (210.25)10 to ()2, ()8, ()16 | June 2006 | 4 |

## Unit-01/Lecture-02

Octal

- Base is 8 and values are from 0 to 7 , largest value is 7 .
- Applications: Avoids large strings of 0 s and 1 s as in Binary system.
- It reduces the size of large binary no.

| 0 | 10 | 20 | 30 | 40 | $\cdots .$. | 70 | 100 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 11 | 21 | 31 | 41 | $\cdots .$. | 71 | 101 |
| 2 | 12 | 22 | 32 | 42 | $\cdots .$. | 72 | 102 |
| 3 | 13 | 23 | 33 | 43 | $\cdots .$. | 73 | 103 |
| 4 | 14 | 24 | 34 | 44 | $\cdots \cdots$. | 74 | 104 |
| .5 | 15 | 25 | 35 | 45 | $\cdots .$. | 75 | 105 |
| 6 | 16 | 26 | 36 | 46 | $\cdots .$. | 76 | 106 |
| 7 | 17 | 27 | 37 | 47 | $\cdots .$. | 77 | 107 |

O-D conversion

$$
\begin{aligned}
372_{8} & =3 \times\left(8^{2}\right)+7 \times\left(8^{1}\right)+2 \times\left(8^{0}\right) \\
& =3 \times 64+7 \times 8+2 \times 1 \\
& =250_{10} \\
24.68 & =2 \times\left(8^{1}\right)+4 \times\left(8^{0}\right)+6 \times\left(8^{-1}\right) \\
& =20.75_{10}
\end{aligned}
$$

D-O conversion

- Separate the integer and fractional parts of the given decimal no.
- Convert the integer part into desired radix
- Convert the fractional part into desired radix
- Combine the results of steps 2 and 3 to get the final answer.
(85.63)10

$(85.63) 10=(1010101.10100) 2$

$$
\begin{aligned}
& \frac{266}{8}=33+\text { remainder of } 2 \mathrm{LSD} \\
& \frac{33}{8}=4+\text { remainder of } 1 \\
& \frac{4}{8}=0+\text { remainder of } 4 \xrightarrow{\text { MSD }} \\
& 266_{10}=412_{8}
\end{aligned}
$$

O-B conversion

| $\stackrel{4}{4}$ | 7 | 2 |
| :---: | :---: | :---: |
| $\downarrow$ | $\downarrow$ | $\downarrow$ |
| 100 | 111 | 010 |

Thus, octal 472 is equivalent to binary 100111010. As another example, consider converting $5431_{8}$ to binary:


Chus, $5431_{8}=101100011001_{2}$.
B-O conversion

- Group binary into 3 digits starting from LSB
- Covert each group into its equivalent decimal. The decimal no is same as octal as each group is restricted to 3 .

$$
\begin{array}{ccccc}
1 & 1 & 0 & 1 & 1_{2} \\
2^{4}+2^{3}+ & 0+2^{1}+2^{0} & =16+8+2+1 \\
& =27_{10}
\end{array}
$$

Let's try another example with a greater number of bits:

$$
\begin{gathered}
1 \\
2^{7}+0+2^{5}+2^{4}+0 \begin{array}{r}
0 \\
0
\end{array}+2^{2}+0+2^{0}=181_{10}
\end{gathered}
$$



## O-H conversion

- Convert the given octal number into equivalent binary
- Then convert this binary into HEX.

Hexadecimal

- Base is 16 , Values are $0,1,2,3,4,5,6,7,8,9, A, B, C, D, E, F$, largest value is $F$
- Applications: HEX no.sare very compact and easy to convert to binary and vice-versa. Used in programming languages.

| 0 | 10 | 20 | 30 | ............. | 90 | A0 | ............. | F0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 11 | 21 | 31 | ............. | 91 | A1 | ............. | F1 |
| 2 | 12 | 22 | 32 | .............. | 92 | A2 | .............. | F2 |
| 3 | 13 | 23 | 33 | .............. | 93 | A3 | ............. | F3 |
| 4 | 14 | 24 | . 34 | .............. | 94 | A4 | .............. | F4 |
| - |  |  | . |  |  |  | [12] |  |
| D | ID | 2D | 3D | ........ | 9D | AD | ............. | FD |
| $E$ | 1E | 2E | 3E | .............. | 9E | AE | ............. | FE |
| F | IF | 2F | 3F | ............. | 9F | AF | ... | EF |

H-Dconversion

$$
\begin{aligned}
356_{16} & =3 \times 16^{2}+5 \times 16^{1}+6 \times 16^{0} \\
& =768+80+6 \\
& =854_{10} \\
2 \mathrm{AF}_{16} & =2 \times 16^{2}+10 \times 16^{1}+15 \times 16^{0} \\
& =512+160+15 \\
& =687_{10}
\end{aligned}
$$

$9 \mathrm{~F} 2_{16}=$

$$
\begin{aligned}
& =100111110010_{2}
\end{aligned}
$$

D-Hconversion
(a) Convert $423_{10}$ to hex.

Solution

$$
\begin{aligned}
& \frac{423}{16}=26+\text { remainder of } 7 \\
& \frac{26}{16}=1+\text { remainder of } 10 \\
& \frac{1}{16}=0+\text { remainder of } 1 \square
\end{aligned} 423_{10}=1 \mathrm{~A} 7_{16}, ~ \$
$$

(b) Convert $214_{10}$ to hex.

Solution

$$
\begin{array}{r}
\frac{214}{16}=13+\text { remainder of } 6 \\
\frac{13}{16}=0+\text { remainder of } 13 \\
214_{10}=\mathrm{D}_{16}
\end{array}
$$

D-Hconversion

Separate the integer and fractional parts :


Convert integer part:

| 18 | 2003 | Hex |  |  |
| ---: | ---: | ---: | :---: | :---: |
| 16 | 125 | 3 |  |  |
| 16 | 7 | $\cdots$ |  |  |$]$| LSD |
| :---: |

$$
\therefore(2003)_{10}=(7 \mathrm{D3})_{16}
$$

Fig. P. 1.9.13(a)
Convert the fractional part into hex:


Fig. P. 9.13 (b)
$\therefore \quad(0.31)_{40}=(0.4 \mathrm{~F} 5 \mathrm{C} 2)_{16}$
Combine the results of steps 2 and 3 :
ombining the results of steps 2 and 3 we get

$$
\therefore \quad(2003,31)_{10}=(7 \mathrm{D} 3.4 \mathrm{~F} 5 \mathrm{C} 2)_{16}
$$

## B-Hconversion

- Break binary into 4 bit sections from LSB to MSB
- Convert each group binary no to its HEX equivalent
$1110100110_{2}=\underbrace{0011}_{3} \underbrace{1010}_{\mathrm{A}} \underbrace{0110}_{6}$

$$
=3 \mathrm{~A} 6_{16}
$$

H-Bconversion

- Convert each HEX digit to its 4 bit binary equivalent
- Combine the 4 bit sections by removing the spaces.


## H-Oconversion

- Represent each HEX digit by a 4 bit binary no.
- Combine these 4 bit binary sections by removing the spaces
- Now group these binary bits into groups of 3 bits starting from LSB side
- Then convert each of this 3 bit group into an octal digit.

| S.NO | RGPV QUESTIONS | Year | Marks |
| :--- | :--- | :--- | :--- |
| Q.1 | Convert (12.0625)10 to Binary | Dec 2006 | 5 |
| Q.2 | Convert (623.77)8 to ()2, ()8, ()16 | Dec 2006 | 5 |
| Q.3 | Convert (2AC5.D)10 to ()2, ()8, ()16 | June 2008 | 5 |

## Unit-03/Lecture-03

## Codes

- No.s, letters or words are represented by a specific group of symbols, it is said that the no is encoded
- The group of symbols is called code
- Codes can be represented in numbers or alphanumeric letters.

Types of Codes


## Weighted Codes

- They obey positional weight principle
- Each position represents specific weight

Non Weighted Codes

- Positional weights are not assigned
- Eg Excess-3 and Gray Codes

BCD
Advanatages:

- Binary equivalents of decimal no.s are only needed to be remembers


## Disadvantages:

- The addition and subtraction has different rules
- BCD arithmetic is bit complicated

D-BCD


As another example, let us change 943 to its BCD-code representation:


BCD-D

Convert 0110100000111001 (BCD) to its decimal equivalent.
Solution
Divide the BCD number into four-bit groups and convert each to decimal.
$\underbrace{0110}_{6} \underbrace{1000}_{8} \underbrace{0011}_{3} \underbrace{1001}_{9}$

BCD-XS3

- Convert BCD to decimal
- Add (3)10 to this decimal number
- Convert into Binary to get the excess 3 code.

BCD Addition

$$
\begin{array}{rrrl}
47 & 0100 & 0111 & \leftarrow \mathrm{BCD} \text { for } 47 \\
+\underline{35} 82 & \frac{0011}{00101} & \leftarrow \mathrm{BCD} \text { for } 35 \\
0111 & 1100 & \leftarrow \text { invalid sum in first digit } \\
1 & \underbrace{0110}_{8} & \leftarrow \underbrace{0}_{2} & \leftarrow \text { add } 6 \text { to correct } \\
0010 & \leftarrow \text { correct BCD sum }
\end{array}
$$

$$
\begin{array}{rr|rl} 
& \begin{array}{c}
\downarrow \\
1
\end{array} & & \\
+\frac{0101}{98} & +\frac{001}{1001} & \leftarrow \operatorname{BCD} \text { for } 59 \\
1000 & \leftarrow \operatorname{BCD} \text { for } 38 \\
1001 & -0001 & \leftarrow \text { perform addition } \\
& \underbrace{1001}_{9} & \underbrace{0111}_{7} & \leftarrow \text { add } 6 \text { to correct } \\
& \text { BCD for } 97
\end{array}
$$

$$
\begin{aligned}
& 2750010 \quad 0111 \quad 0101 \leftarrow \mathrm{BCD} \text { for } 275 \\
& +\underline{641}+\underline{011001000001 \leftarrow \operatorname{BCD} \text { for } 641} \\
& 9161000 \quad 1011 \quad 0110 \leftarrow \text { perform addition } \\
& +\frac{0110}{}+0.4 \text { add } 6 \text { to correct second digit }
\end{aligned}
$$

## Excess 3

- It is no weighted code
- The excess 3 code word are derived from the 8421 BCD code word by adding (0011)2 or (3)10 to each word in 8421.
Decimal Number $\rightarrow 8421 \mathrm{BCD} \underset{0011}{\text { Add }}$ Excess -3 code

| Decimal | BCD |  |  |  |  | Excess - 3 |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 8 | 4 | 2 | 1 | BCD +00 |  |  | 1 | 1 |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 |  |
| 1 | 0 | 0 | 0 | 1 | 0 | 1 | 0 | 0 |  |
| 2 | 0 | 0 | 1 | 0 | 0 | 1 | 0 | 1 |  |
| 3 | 0 | 0 | 1 | 1 | 0 | 1 | 1 | 0 |  |
| 4 | 0 | 1 | 0 | 0 | 0 | 1 | 1 | 1 |  |
| 5 | 0 | 1 | 0 | 1 | 1 | 0 | 0 | 0 |  |
| 6 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 1 |  |
| 7 | 0 | 1 | 1 | 1 | 1 | 0 | 1 | 0 |  |
| 8 | 1 | 0 | 0 | 0 | 1 | 0 | 1 | 1 |  |
| 9 | 1 | 0 | 0 | 1 | 1 | 1 | 0 | 0 |  |


| Decimal | 1 Binary | Octal | Hexadecimal | BCD | GRAY |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0000 | 0000 |  |
| 1 | 1 | 1 | 1 | 0001 | 0001 |  |
| 2 | 10 | 2 | 2 | 0010 | 0011 |  |
| 3 | 11 | 3 | 3 | 0011 | 0010 |  |
| 4 | 100 | 4 | 4 | 0100 | 0110 |  |
| 5 | 101 | 5 | 5 | 0101 | 0111 |  |
| 6 | 110 | 6 | 6 | 0110 | 0101 |  |
| 7 | 111 | 7 | 7 | 0111 | 0100 |  |
| 8 | 1000 | 10 | 8 | 1000 | 1100 |  |
| 9 | 1001 | 11 | 9 | 1001 | 1101 |  |
| 10 | 1010 | 12 | A | 00010000 | 1111 |  |
| 11 | 1011 | 13 | B | 00010001 | 1110 |  |
| 12 | 1100 | 14 | c | 00010010 | 1010 |  |
| 13 | 1101 | 15 | D | 00010011 | 1011 |  |
| 14 | 1110 | 16 | E | 00010100 | 1001 |  |
| 15 | 1111 | 17 | F | 00010101 | 1000 |  |
| S.NO | RGPV QUESTIONS |  |  |  | Year | Marks |
| Q. 1 | Convert (101111101)2 to ()8, (C346)10 to ()2, (3906)10 to ()BCD, (370)8 to ()16 |  |  |  | Jun 2005 | 10 |
| Q. 2 | Convert (428)10 | () XS3 |  |  | Dec 2010 | 5 |

## Unit-01/Lecture-04

## GRAY

- It is non weighted code. No specific weights are assigned to the bit positions
- Only one bit change each time the decimal number is incremented called unit distance code
- Applications: used in Shaft position encoders (both linear and angular)

| $B_{2}$ | $B_{1}$ | $B_{0}$ | $G_{2}$ | $G_{1}$ | $G_{0}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 | 1 |
| 0 | 1 | 0 | 0 | 1 | 1 |
| 0 | 1 | 1 | 0 | 1 | 0 |
| 1 | 0 | 0 | 1 | 1 | 0 |
| 1 | 0 | 1 | 1 | 1 | 1 |
| 1 | 1 | 0 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 | 0 | 0 |


(a)

(b)

## G-B

- The MSB of Gray and Binary are the same. So write it directly
- Add binary MSB to the next bit of Gray code. Record the result and ignore the carries
- Continue this process until the LSB is reached.



## B-G

- Record the MSB as it is.
- Add this bit to the next position, recording the sum and neglecting any carry
- Record successive sums until completed.

Convert binary 1011 to gray.

Given Binary number: MSB 1 LSB

Gray code : MSB LSB


MSB directly

$$
\therefore(1011)_{2}=(1110)_{\text {gray }}
$$

Aphanumeric Codes

- Represent numbers as well as alphabetic characters
- Some represent symbols and instructions as well.
- Eg ASCII, EBCDIC (Extended Binary Coded Decimal Interchange Code), Hollerith

ASCII Codes (American Standard Code for Information Interchange)

- ASCII has 128 characters and symbols and requires 7 bits
- It can be considered as 8 bit with $\mathrm{MSB}=0$
- This 8 bit is 00 to 7 F
- It has 32 control symbols those are not displayed on screen
- 94 printable characters and SPACE and DEL

| Character | HEX | Decimal | Character | HEX | Decimal | Character | HEX | Decimal | Character | HEX | Decimal |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| NUL (null) | 0 | 0 | Space | 20 | 32 | @ | 40 | 64 | . | 60 | 96 |
| Start Heading | 1 | 1 | 1 | 21 | 33 | A | 41 | 65 | a | 61 | 97 |
| Start Text | 2 | 2 | " | 22 | 34 | B | 42 | 66 | b | 62 | 98 |
| End Text | 3 | 3 | \# | 23 | 35 | C | 43 | 67 | c | 63 | 99 |
| End Transmit. | 4 | 4 | \$ | 24 | 36 | D | 44 | 68 | d | 64 | 100 |
| Enquiry | 5 | 5 | \% | 25 | 37 | E | 45 | 69 | e | 65 | 101 |
| Acknowlege | 6 | 6 | \& | 26 | 38 | F | 46 | 70 | f | 66 | 102 |
| Bell | 7 | 7 | , | 27 | 39 | G | 47 | 71 | g | 67 | 103 |
| Backspace | 8 | 8 | 1 | 28 | 40 | H | 48 | 72 | h | 68 | 104 |
| Horiz. Tab | 9 | 9 | ) | 29 | 41 | 1 | 49 | 73 | i | 69 | 105 |
| Line Feed | A | 10 | - | 2 A | 42 | $J$ | 4 A | 74 | j | 6 A | 106 |
| Vert. Tab | B | 11 | + | 2 B | 43 | K | 4B | 75 | k | 6 B | 107 |
| Form Feed | c | 12 | , | 2 C | 44 | L | 4 C | 76 | I | 6 C | 108 |
| Carriage Return | D | 13 | - | 2D | 45 | M | 4D | 77 | m | 6 D | 109 |
| Shift Out | E | 14 | . | 2 E | 46 | N | 4 E | 78 | n | 6 E | 110 |
| Shift in | F | 15 | 1 | 2 F | 47 | 0 | 4F | 79 | - | 6 F | 111 |
| Data Link Esc | 10 | 16 | 0 | 30 | 48 | P | 50 | 80 | $p$ | 70 | 112 |
| Direct Control 1 | 11 | 17 | 1 | 31 | 49 | Q | 51 | 81 | q | 71 | 113 |
| Direct Control 2 | 12 | 18 | 2 | 32 | 50 | R | 52 | 82 | r | 72 | 114 |
| Direct Control 3 | 13 | 19 | 3 | 33 | 51 | S | 53 | 83 | s | 73 | 115 |
| Direct Control 4 | 14 | 20 | 4 | 34 | 52 | T | 54 | 84 | t | 74 | 116 |
| Negative ACK | 15 | 21 | 5 | 35 | 53 | U | 55 | 85 | $u$ | 75 | 117 |
| Synch Idle | 16 | 22 | 6 | 36 | 54 | v | 56 | 86 | $v$ | 76 | 118 |
| End Trans Block | 17 | 23 | 7 | 37 | 55 | w | 57 | 87 | w | 77 | 119 |
| Cancel | 18 | 24 | 8 | 38 | 56 | X | 58 | 88 | x | 78 | 120 |
| End of Medium | 19 | 25 | 9 | 39 | 57 | Y | 59 | 89 | y | 79 | 121 |
| Substitue | 1A | 26 | : | 3 A | 58 | z | 5 A | 90 | z | 7A | 122 |
| Escape | 1 B | 27 | ; | 3 B | 59 | I | 5 B | 91 | ( | 78 | 123 |
| Form separator | 1 C | 28 | < | 3 C | 60 | 1 | 5 C | 92 | 1 | 7 C | 124 |
| Group separator | 1 D | 29 | $=$ | 3 D | 61 | 1 | 50 | 93 | ) | 7 D | 125 |
| Record Separator | 1E | 30 | > | 3 E | 62 | $\wedge$ | 5E | 94 | ~ | 7 E | 126 |
| Unit Separator | 1F | 31 | ? | 3 F | 63 | - | 5 F | 95 | Delete | 7F | 127 |

## Unit-01/Lecture-05

## Theorems and Properties

## BOOLEAN THEOREMS

1. $x \cdot 0=0$
2. $x+1=x$
3. $x-\bar{x}=0$
4. $x+0=x$
5. $x+x=x$
6. $x+\bar{x}=1$
7. $x+y=y \cdot x$
8. $x+(y+z)=(x+y)+z=x+y+z$
9. $x-x=x$
13b. $(w+x)(y+z)=w y+x y+w z+x z$
10. $x+x y=x$
3a. $\quad x(y+z)=x y+x z$
15b. $\bar{x}+x y=\bar{x}+y$
11. $\overline{x+y}=\bar{x} \bar{y}$
5a. $x+\bar{x} y=x+y$
12. $x+1=1$
13. $x+y=y+x$
14. $x(y z)=(x y) z=x y z$
15. $\overline{x y}=\vec{x}+\vec{y}$

Boolean Functions

| 1. | Commutative Law | $\mathrm{A} \cdot \mathrm{B}=\mathrm{B} \cdot \mathrm{A}$ <br> $\mathrm{A}+\mathrm{B}=\mathrm{B}+\mathrm{A}$ |
| :--- | :--- | :--- |
| 2. | Associative Law | $\mathrm{A} \cdot \mathrm{B}) \cdot \mathrm{C}=\mathrm{A} \cdot(\mathrm{B} \cdot \mathrm{C})$ <br> $(\mathrm{A}+\mathrm{B})+\mathrm{C}=\mathrm{A}+(\mathrm{B}+\mathrm{C})$ |
| 3. | Distributive Law | $\mathrm{A} \cdot(\mathrm{B}+\mathrm{C})=\mathrm{AB}+\mathrm{AC}$ |
| 4. | AND Laws | $\mathrm{A} \cdot 0=0$ |
|  |  | $\mathrm{~A} \cdot \mathrm{I}=\mathrm{A}$ |
| $\mathrm{A} \cdot \mathrm{A}=\mathrm{A}$ |  |  |, | $\mathrm{A} \cdot \overline{\mathrm{A}}=0$ |
| :--- |

```
To prove that \(\mathbf{A}+\mathbf{A B}=\mathbf{A}\) :
\[
\begin{array}{rlrl}
\text { LHS } & =A+A B=A(1+B) \\
\text { But } & =1+B) & =1 \quad \therefore \text { LHS }=\mathrm{A} \cdot 1 \quad \text { But } \mathrm{A} \cdot 1=\mathrm{A} \\
\therefore \quad \mathrm{LHS} & =\mathrm{A}=\mathrm{RHS} \\
\therefore \quad \mathrm{~A}+\mathrm{AB} & =\mathrm{A}
\end{array}
\]
```

To prove that $\mathbf{A}+\overline{\mathbf{A}} \mathbf{B}=\mathbf{A}+\mathbf{B}$ :

$$
\begin{aligned}
& \text { LHS }=A+\bar{A} B=A+A B+\bar{A} B \\
\therefore \quad & \text { LHS }=A+B(A+\bar{A}) \\
=1 & \therefore \quad \text { LHS }=A+(B \cdot 1)=A+B \\
& \therefore \quad \text { LHS }=\text { RHS } A=A+A B
\end{aligned}
$$

it $(\mathrm{A}+\overline{\mathrm{A}})=1$

$$
\therefore \mathrm{A}+\overline{\mathrm{A}} \mathrm{~B}=\mathrm{A}+\mathrm{B}
$$

To prove that $(\mathbf{A}+\mathbf{B})(\mathbf{A}+\mathbf{C})=\mathbf{A}+\mathbf{B C}$ :

$$
\mathrm{LHS}=(\mathrm{A}+\mathrm{B})(\mathrm{A}+\mathrm{C})
$$

$\therefore$ LHS $=\mathrm{AA}+\mathrm{AC}+\mathrm{BA}+\mathrm{BC} \quad \ldots . .$. According to the distributive law.
But $A A=A, \quad \therefore$ LHS $=A+A C+B A+B C$
$\therefore$ LHS $=A(1+C)+A B+B C$
But $(1+C)=1$
$\therefore$ LHS $=A+A B+B C=A(1+B)+B C$
But $(1+B)=1$

## Unit-01/Lecture-06

Canonical Forms

- The meaning of Canonical is conforming to a general rule
- This rule states that each term used in a switching equation must contain all the available input variables.
- It can be expressed in SOP or POS forms
- When we simplify sometimes input variables are eliminated
- Canonical expressions are not simplified hence contains redundancies.
- These equations are not in the minimized form because each term contains all the literals

Standard Form

- Each term can contain 1,2 or any no of literals.
- Can be expressed in SOP or POS forms.

| 1. | $\mathrm{Y}=\mathrm{AB}+\mathrm{AB} \overline{\mathrm{C}}+\overline{\mathrm{A} B C}$ | Standard SOP |
| :--- | :--- | :--- |
| 2. | $\mathrm{Y}=\mathrm{AB}+\mathrm{A} \overline{\mathrm{B}}+\overline{\mathrm{A}} \overline{\mathrm{B}}$ | Canonical SOP |
| 3. | $\mathrm{Y}=(\overline{\mathrm{A}}+\mathrm{B}) \cdot(\mathrm{A}+\mathrm{B}) \cdot(\mathrm{A}+\overline{\mathrm{B}})$ | Canonical POS |
| 4. | $\mathrm{Y}=(\overline{\mathrm{A}}+\mathrm{B}) \cdot(\mathrm{A}+\mathrm{B}+\mathrm{C})$ | Standard POS |

1. 

$$
\begin{aligned}
& \mathrm{Y}=\underbrace{\mathrm{ABC}}_{\mathrm{m}_{7}}+\underbrace{\bar{A} B C}_{\mathrm{m}_{3}}+\underbrace{\mathrm{AB} \overline{\mathrm{~B}} \overline{\mathrm{C}}}_{\mathrm{m}_{4}} \leftarrow \text { Given logic expression } \\
& \hline
\end{aligned}
$$

where $\Sigma$ denotes sum of products.
2.

$$
\begin{aligned}
\mathbf{Y} & =(\underbrace{\mathrm{A}+\overline{\mathrm{B}}+\mathrm{C}}_{\mathbf{M}_{2}})(\underbrace{\mathrm{A}+\mathrm{B}+\mathrm{C}}_{\mathbf{M}_{0}})
\end{aligned} \underbrace{(\underbrace{\bar{A}+\overline{\mathrm{B}}+\mathrm{C}})}_{\mathbf{M}_{6}} \leftarrow \text { Corresponding ma } \text { Given expressic }
$$

where $\Pi$ denotes product of sums.

| Variables |  |  | Minterms | Maxterms |
| :---: | :---: | :---: | :---: | :---: |
| A | B | C | $\mathrm{m}_{1}$ | $\mathrm{M}_{1}$ |
| 0 | 0 | 0 | $\overline{\mathrm{~A}} \overline{\mathrm{~B}} \overline{\mathrm{C}}=\mathrm{m}_{0}$ | $\mathrm{~A}+\mathrm{B}+\mathrm{C}=\mathrm{M}_{0}$ |
| 0 | 0 | 1 | $\overline{\mathrm{~A}} \overline{\mathrm{~B}} \mathrm{C}=\mathrm{m}_{1}$ | $\mathrm{~A}+\mathrm{B}+\overline{\mathrm{C}}=\mathrm{M}_{1}$ |
| 0 | 1 | 0 | $\overline{\mathrm{~A}} \mathrm{~B} \overline{\mathrm{C}}=\mathrm{m}_{2}$ | $\mathrm{~A}+\overline{\mathrm{B}}+\mathrm{C}=\mathrm{M}_{2}$ |
| 0 | 1 | 1 | $\overline{\mathrm{~A}} \mathrm{BC}=\mathrm{m}_{3}$ | $\mathrm{~A}+\overline{\mathrm{B}}+\overline{\mathrm{C}}=\mathrm{M}_{3}$ |
| 1 | 0 | 0 | $\mathrm{~A} \overline{\mathrm{~B}} \overline{\mathrm{C}}=\mathrm{m}_{4}$ | $\overline{\mathrm{~A}}+\mathrm{B}+\mathrm{C}=\mathrm{M}_{4}$ |
| 1 | 0 | 1 | $\mathrm{~A} \overline{\mathrm{~B}} \mathrm{C}=\mathrm{m}_{5}$ | $\overline{\mathrm{~A}}+\mathrm{B}+\overline{\mathrm{C}}=\mathrm{M}_{5}$ |
| 1 | 1 | 0 | $\mathrm{~A} B \overline{\mathrm{C}}=\mathrm{m}_{6}$ | $\overline{\mathrm{~A}}+\overline{\mathrm{B}}+\mathrm{C}=\mathrm{M}_{6}$ |
| 1 | 1 | 1 | $\mathrm{~A} \mathrm{BC}=\mathrm{m}_{7}$ | $\overline{\mathrm{~A}}+\overline{\mathrm{B}}+\overline{\mathrm{C}}=\mathrm{M}_{7}$ |

Ex: Express $f(A, B, C)=\left(A^{\prime}+B\right)\left(B^{\prime}+C\right)$ in sum of min terms and product of max terms Dec 06, $5 M$
$\mathrm{f}(\mathrm{A}, \mathrm{B}, \mathrm{C})=(\overline{\mathrm{A}}+\mathrm{B})(\overline{\mathrm{B}}+\mathrm{C})$
Simplify the expression to get canonical POS form :

$$
\begin{aligned}
\mathrm{f}(\mathrm{~A}, \mathrm{~B}, \mathrm{C}) & =(\overline{\mathrm{A}}+\mathrm{B}+\mathrm{C} \overline{\mathrm{C}})(\overline{\mathrm{B}}+\mathrm{C}+\mathrm{A} \overline{\mathrm{~A}}) \\
& =(\overline{\mathrm{A}}+\mathrm{B}+\mathrm{C})(\overline{\mathrm{A}}+\mathrm{B}+\overline{\mathrm{C}})(\mathrm{A}+\overline{\mathrm{B}}+\mathrm{C})(\overline{\mathrm{A}}+\overline{\mathrm{B}}+\mathrm{C}) \\
& =\mathrm{M}_{4} \mathrm{M}_{5} \mathrm{M}_{2} \mathrm{M}_{6} \\
\therefore \mathrm{f}(\mathrm{~A}, \mathrm{~B}, \mathrm{C}) & =\Pi \mathrm{M}(2,4,5,6)
\end{aligned}
$$

Simplify the expression to get canonical SOP form :

$$
\begin{aligned}
\mathrm{f}(\mathrm{~A}, \mathrm{~B}, \mathrm{C}) & =(\overline{\mathrm{A}}+\mathrm{B})(\overline{\mathrm{B}}+\mathrm{C})=\overline{\mathrm{A}} \overline{\mathrm{~B}}+\overline{\mathrm{A}} \mathrm{C}+\mathrm{B} \overline{\mathrm{~B}}+\mathrm{BC} \\
& =\overline{\mathrm{A}} \overline{\mathrm{~B}}+\overline{\mathrm{A} C}+\mathrm{BC}(\because \mathrm{~A} \overline{\mathrm{~A}}=0) \\
& =\overline{\mathrm{A}} \overline{\mathrm{~B}}(\mathrm{C}+\overline{\mathrm{C}})+\overline{\mathrm{A} C} \mathrm{C}(\mathrm{~B}+\overline{\mathrm{B}})+\mathrm{BC}(\mathrm{~A}+\overline{\mathrm{A}}) \\
& =\overline{\mathrm{A}} \overline{\mathrm{~B} C}+\overline{\mathrm{A}} \overline{\mathrm{~B}} \overline{\mathrm{C}}+\overline{\mathrm{A}} \mathrm{BC}+\overline{\mathrm{A}} \overline{\mathrm{~B}} \mathrm{C}+\mathrm{ABC}+\overline{\mathrm{A}} \mathrm{BC} \\
& =\overline{\mathrm{A}} \overline{\mathrm{~B}} \mathrm{C}+\overline{\mathrm{A} \overline{\mathrm{~B}} \overline{\mathrm{C}}+\overline{\mathrm{A}} \mathrm{BC}+\mathrm{ABC}=\mathrm{m}_{1}+\mathrm{m}_{3}+\mathrm{m}_{3}+\mathrm{m}_{7}} \\
\mathrm{f}(\mathrm{~A}, \mathrm{~B}, \mathrm{C}) & =\Sigma \mathrm{m}(0,1,3,7)
\end{aligned}
$$

De Morgans Theorem

$$
\begin{aligned}
& (\overline{x+y})=\bar{x} \cdot \bar{y} \\
& (\overline{x \cdot y})=\bar{x}+\bar{y} \\
& \overline{x+y+z}=\bar{x} \cdot \bar{y} \cdot \bar{z} \\
& \overline{x \cdot y \cdot z}=\bar{x}+\bar{y}+\bar{z}
\end{aligned}
$$

$$
\begin{aligned}
& \text { Example } 1 \\
& z=\overline{A+\bar{B} \cdot C} \\
& \text { Example } 2 \\
& =\bar{A} \cdot(\overline{\bar{B} \cdot C}) \\
& \omega=\overline{(A+B C) \cdot(D+E F)} \\
& =\bar{A} \cdot(\overline{\bar{B}}+\bar{C}) \\
& =(\overline{A+B C})+(\overline{D+E F}) \\
& =\bar{A} \cdot(B+\bar{C}) \\
& =(\bar{A} \cdot \overline{B C})+(\bar{D} \cdot \overline{E F}) \\
& =[\bar{A} \cdot(\bar{B}+\bar{C})]+[\bar{D} \cdot(\bar{E}+\bar{F})] \\
& =\bar{A} \bar{B}+\bar{A} \bar{C}+\bar{D} \bar{E}+\bar{D} \bar{F} \\
& x=\overline{\overline{A B}} \cdot \overline{C D} \cdot \overline{E F} \\
& =\overline{\overline{A B}}+\overline{\overline{C D}}+\overline{\overline{E F}} \\
& =A B+C D+E F
\end{aligned}
$$

## Digital Logic Gates

## Truth Table:

Consists of all the possible combinations of the inputs and the corresponding state of output of a logic gate.

## LOGIC GATE TRUTH TABLES

|  |  | OR | NOR | AND | NAND | XOR | XNOR |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A | B | $\mathrm{A}+\mathrm{B}$ | $\frac{\mathrm{A}+\mathrm{B}}{}$ | $\mathrm{A} \cdot \mathrm{B}$ | $\frac{\mathrm{A}-\mathrm{B}}{}$ | $\mathrm{A} \mathrm{\oplus B}$ | $\frac{\mathrm{~A} \oplus \mathrm{~B}}{}$ |
| 0 | 0 | 0 | 1 | 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | 0 | 0 | 1 | 1 | 0 |
| 1 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |
| 1 | 1 | 1 | 0 | 1 | 0 | 0 | 1 |

Basic Gates: NOT, AND, OR
Universal Gates: NAND, NOR

Derived Gates: EX-OR, EX-NOR





Quads

(a)

(b)

(c)

|  | $\bar{C} \bar{D}$ | $\overline{C D}$ | $C D$ | $C \bar{D}$ |
| :--- | :--- | :--- | :--- | :--- |
| $\bar{A} \bar{B}$ | 0 | 0 | 0 | 0 |
| $\bar{A} B$ | 0 | 0 | 0 | 0 |
| $A B$ | 1 | 0 | 0 | 1 |
| $A \bar{B}$ | 1 | 0 | 0 | 1 |

(d)

(e)

Octets

(a)

(c)

(b)

(d)

## Steps for K Map Reduction

1. Place 1s according to the Truth Table or logical expression and fill Os at all the other positions
2. Locate the isolated 1s (That can not be combined) and circle them.
3. Locate 1 s which form a double and encircle them in a group
4. Locate 1 s which form a Quad and encircle them in a group
5. Locate 1 s which form a Octet and encircle them in a group
6. After identifying, check if any 1 s is yet to be encircled, if yes encircle them with each other or with the already encircled 1 s by means of overlapping.
7. Note that the number of groups should be minimum
8. That any 1 can be included any number of times without affecting the expression
9. The redundant group has to be eliminated as it increases the number of gates required.
K Map Examples


$$
\begin{aligned}
& X=\underbrace{\overline{A B}}_{\text {loop 5, }}+\underbrace{\mathrm{BC}}_{\text {loop 5 },}+\underbrace{\overline{\mathrm{ACD}}}_{\text {loop }} \\
& 6.7,8 \quad 6,9,10 \quad 3,7
\end{aligned}
$$

(b)

$X=\underbrace{A B \bar{C}}_{9,10}+\underbrace{\overline{\mathrm{A} \overline{C D}}}_{2,6}+\underbrace{\overline{\mathrm{ABC}}}_{7,8}+\underbrace{\mathrm{ACD}}_{11,15}$
(c)

